

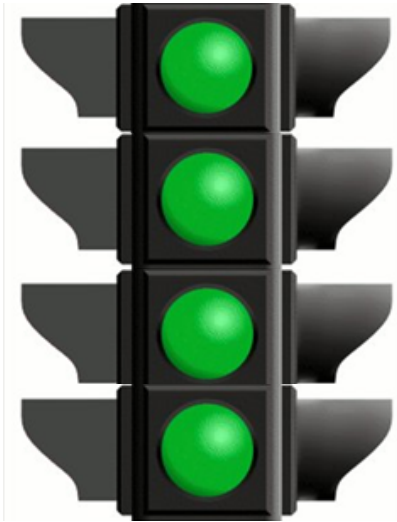


# Segment Routing

Berényi Áron

# Deployment

- In CY2015, SR will be deployed in all of these markets



WEB

SP Core/Edge

SP Agg/Metro

Large Enterprise

# Agenda

- Technology Overview
- Use Cases
- Control and Data Plane
- Traffic Protection



# Technology Overview

# Segment Routing

- Source Routing

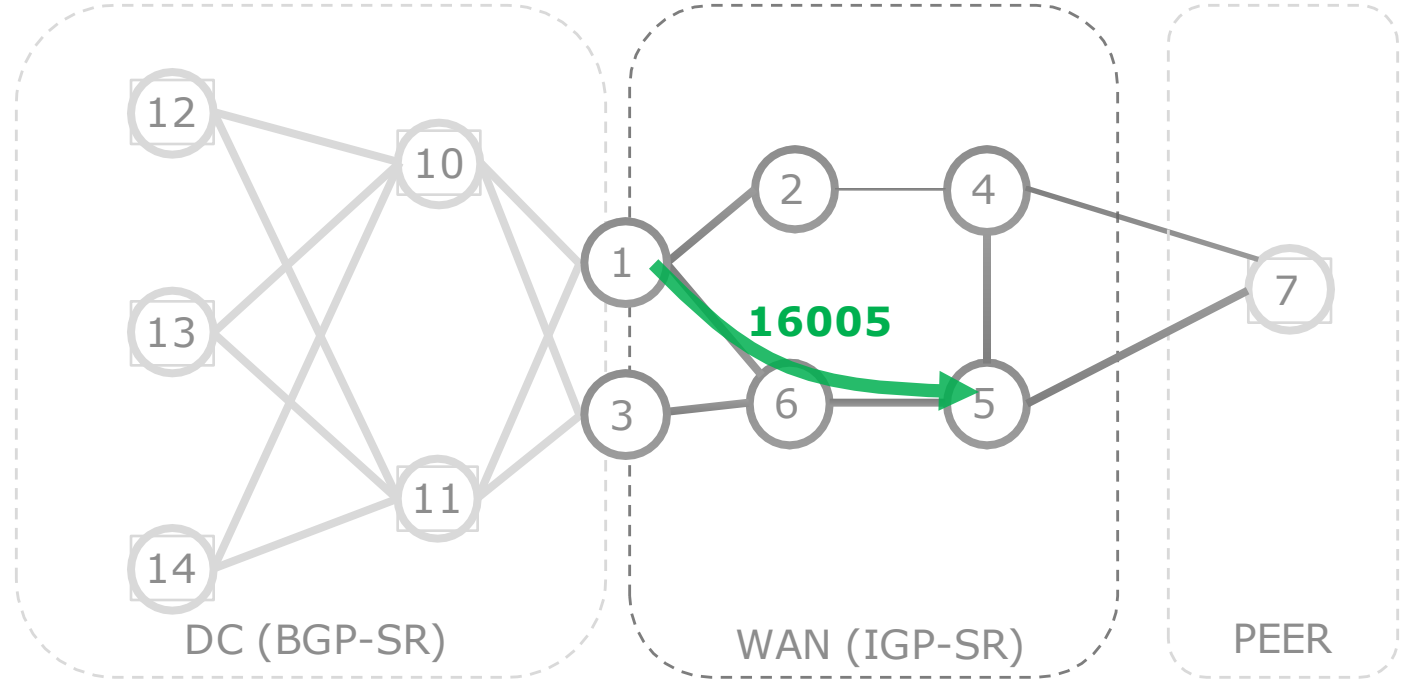
- the source chooses a path and encodes it in the packet header as an ordered list of segments
- the rest of the network executes the encoded instructions without any further per-flow state

- Segment: an identifier for any type of instruction

- forwarding or service

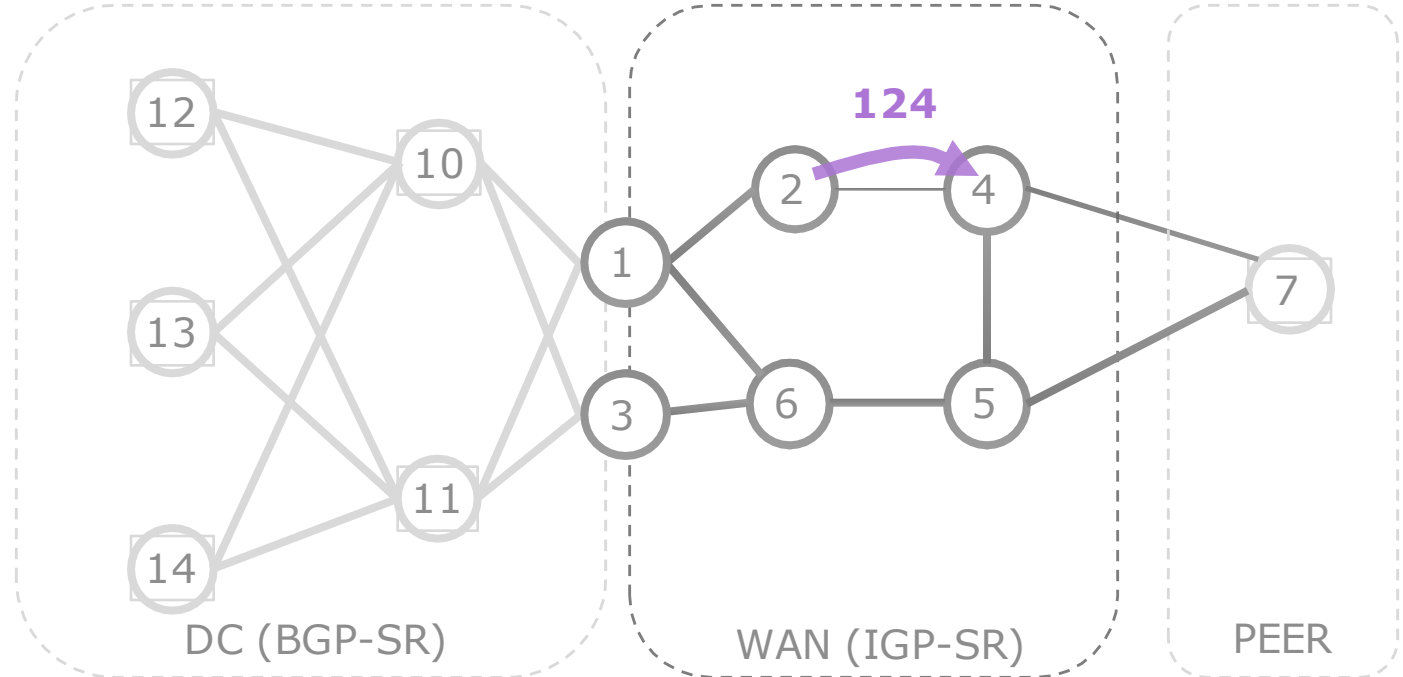
# IGP Prefix Segment

- Shortest-path to the IGP prefix
- Global
- 16000 + Index
- Signaled by ISIS/OSPF



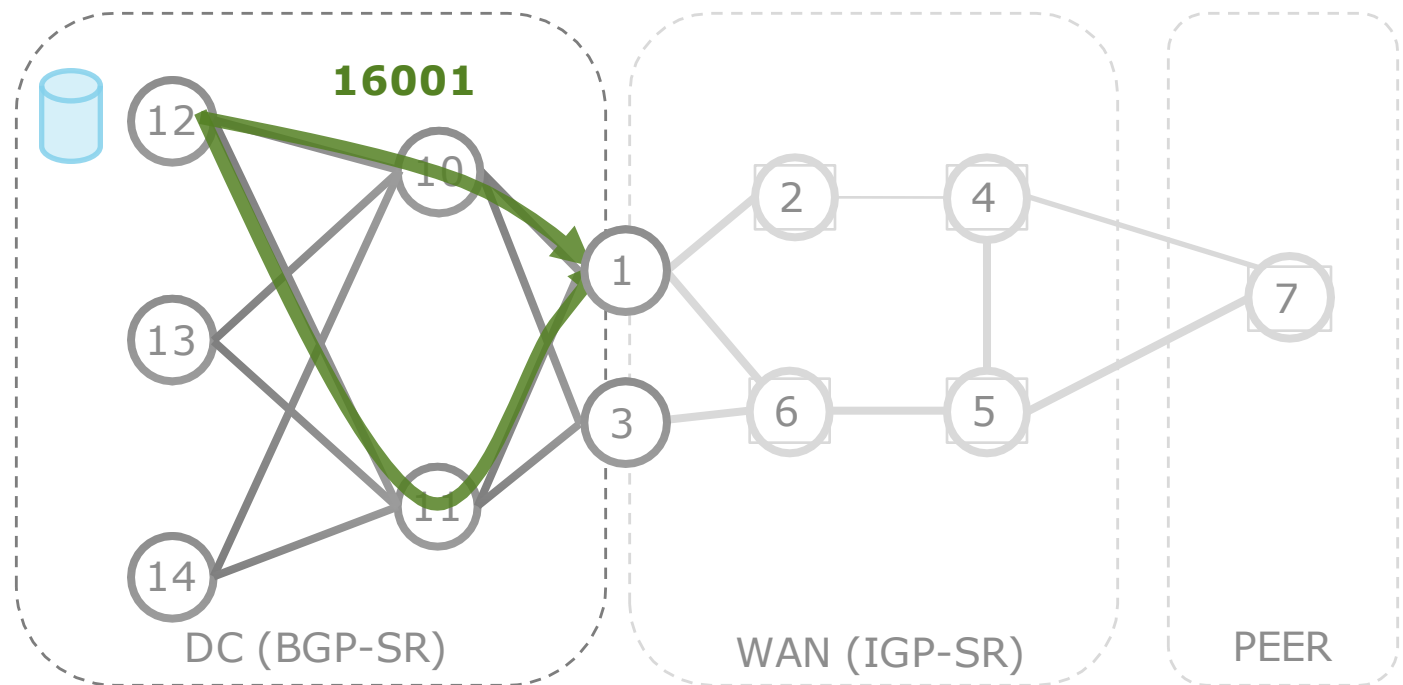
# IGP Adjacency Segment

- Forward on the IGP adjacency
- Local
- 1XY
  - X is the “from”
  - Y is the “to”
- Signaled by ISIS/OSPF



# BGP Prefix Segment

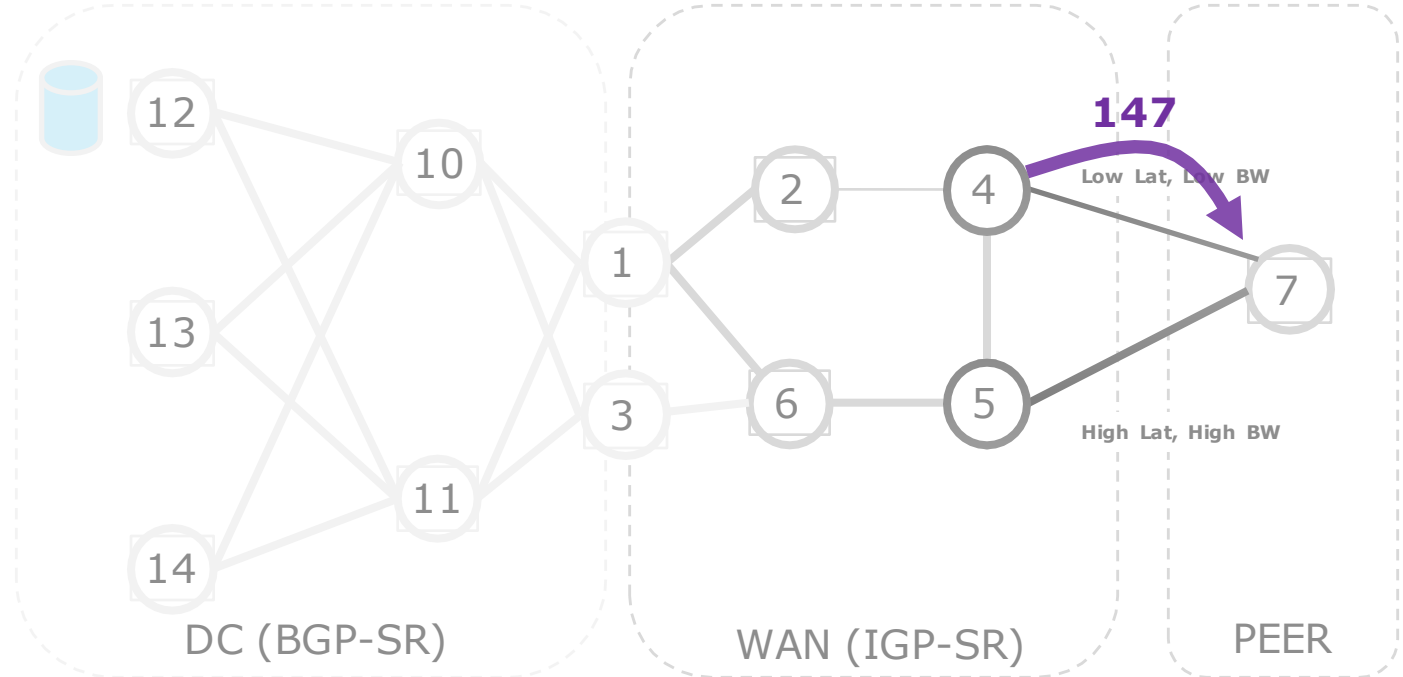
- Shortest-path to the BGP prefix
- Global
- 16000 + Index
- Signaled by BGP





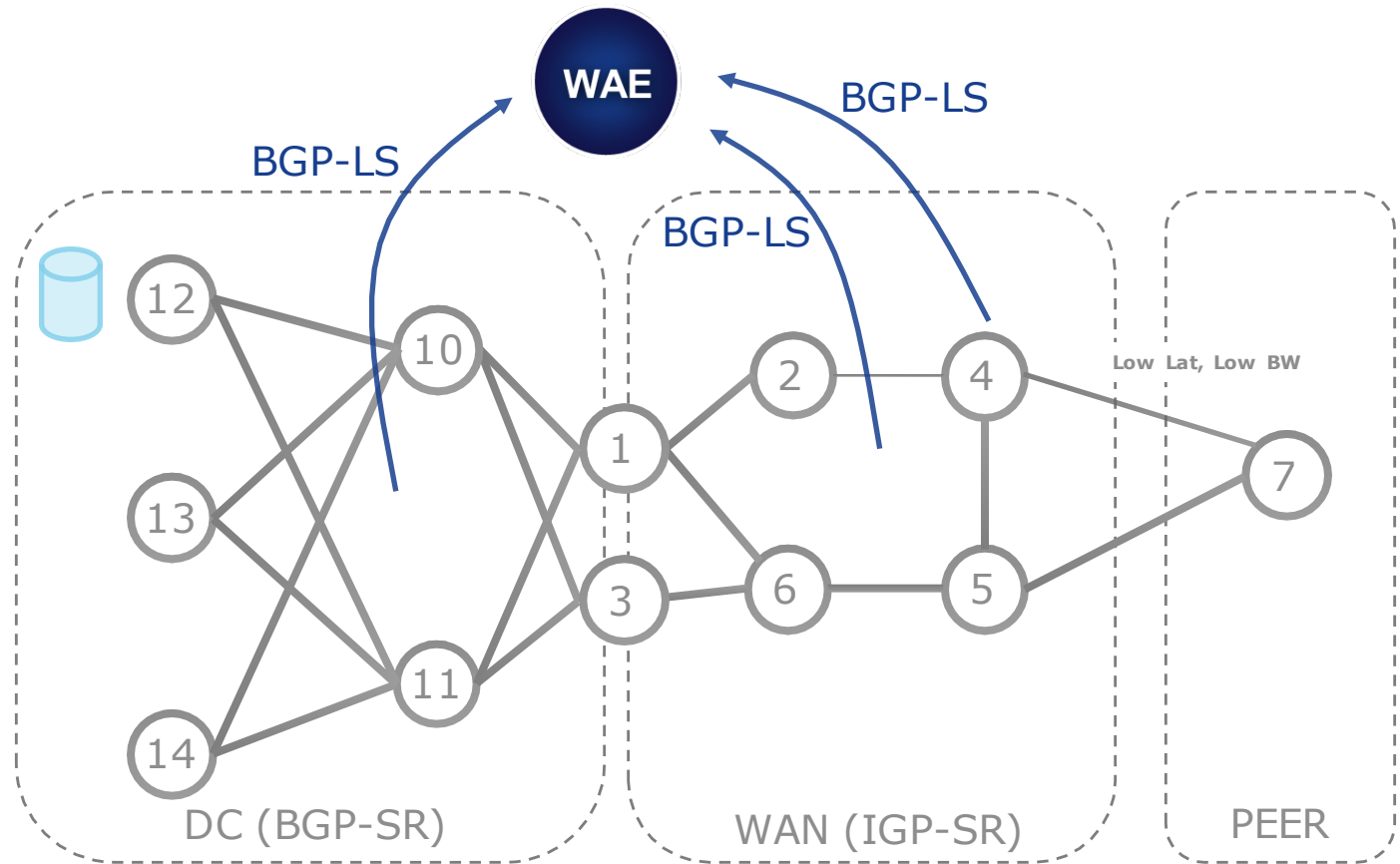
# BGP Peering Segment

- Forward to the BGP peer
- Local
- 1XY
  - X is the “from”
  - Y is the “to”
- Signaled by BGP-LS (topology information) to the controller



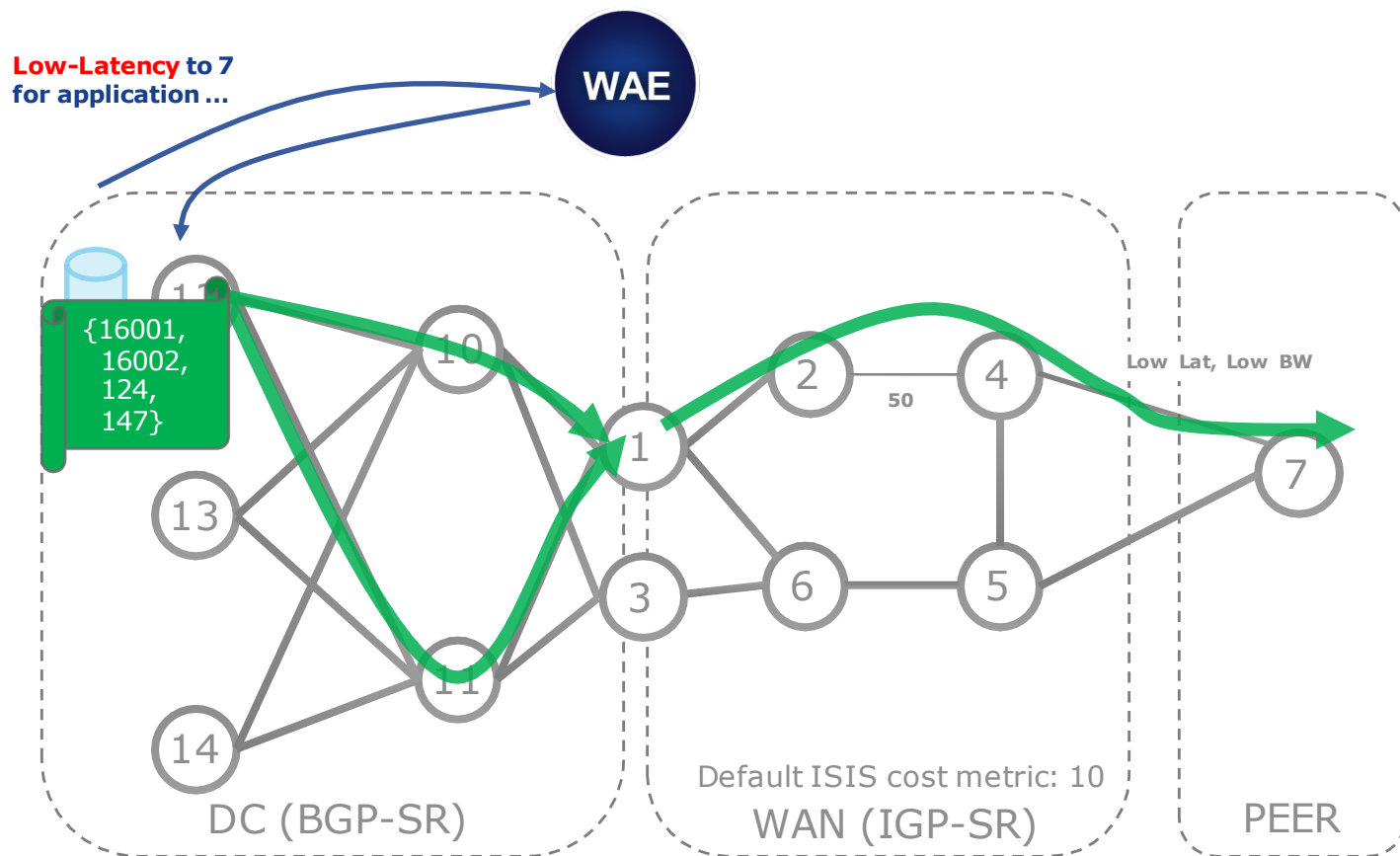
# WAN Automation Engine

- WAE collects via BGP-LS
  - IGP segments
  - BGP segments
  - Topology



# An end-to-end path as a list of segments

- WAE computes that the green path can be encoded as
  - 16001
  - 16002
  - 124
  - 147
- WAE programs a single per-flow state to create an application-engineered end-to-end policy



# Segment Routing Standardization

- IETF standardization in SPRING working group
- Protocol extensions progressing in multiple groups
  - IS-IS
  - OSPF
  - PCE
  - IDR
  - 6MAN
- Broad vendor and customer support

Sample IETF Documents
Segment Routing Architecture ( <a href="#">draft-ietf-spring-segment-routing</a> )
Problem Statement and Requirements ( <a href="#">draft-ietf-spring-problem-statement</a> )
IPv6 SPRING Use Cases ( <a href="#">draft-ietf-spring-ipv6-use-cases</a> )
Segment Routing Use Cases ( <a href="#">draft-filsfils-spring-segment-routing-use-cases</a> )
Topology Independent Fast Reroute using Segment Routing ( <a href="#">draft-francois-spring-segment-routing-ti-lfa</a> )
IS-IS Extensions for Segment Routing ( <a href="#">draft-ietf-isis-segment-routing-extensions</a> )
OSPF Extensions for Segment Routing ( <a href="#">draft-ietf-ospf-segment-routing-extensions</a> )
PCEP Extensions for Segment Routing ( <a href="#">draft-ietf-pce-segment-routing</a> )

Close to 30 IETF drafts in progress

# Segment Routing Product Support


- Platforms: ASR9000, CRS-1/CRS-3, ASR1000, ASR9XX, ISR4XXX
- IS-IS IPv4/IPv6
  - Node/Adjacency SID advertisement
  - LDP interworking (mapping server/client)
  - Traffic protection (Topology Independent LFA link protection)
- OSPFv2
  - Node SID advertisement
  - Traffic Protection (LFA)
- SR Traffic Engineering manual/PCEP
- OAM ping/trace

# Application Engineered Routing Journey

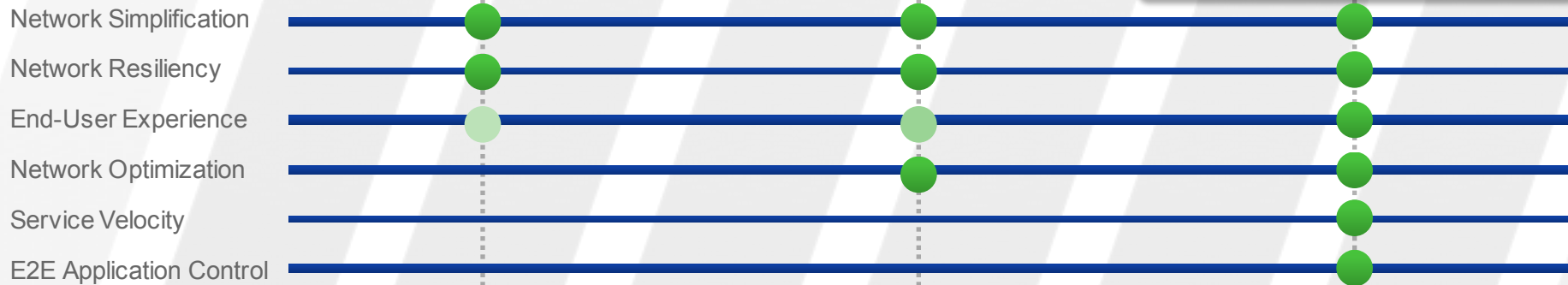
Adding value at your own pace

 Enable Segment Routing on the network (Software only)

 Insert Orchestration, SDN controller

 Connect with Cisco's and third party VNFs

## Benefits

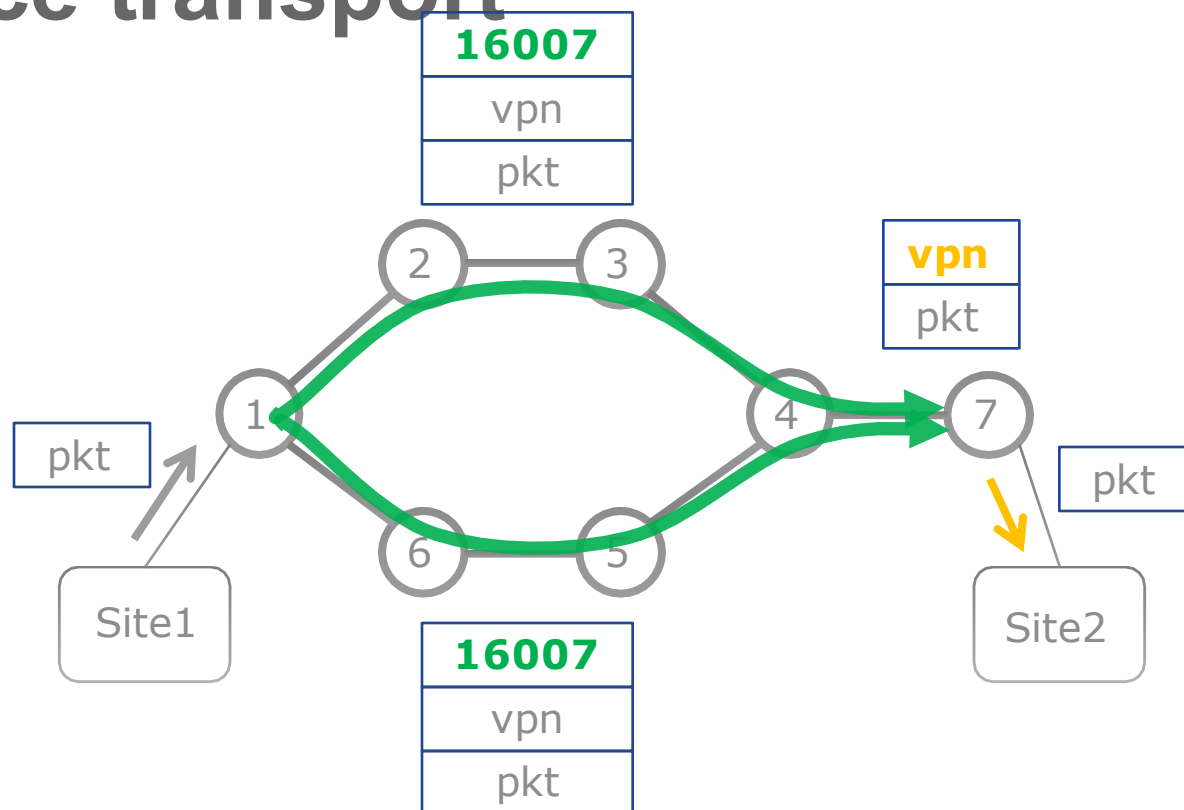




# Use-Cases

# IPv4/6 VPN/Service transport

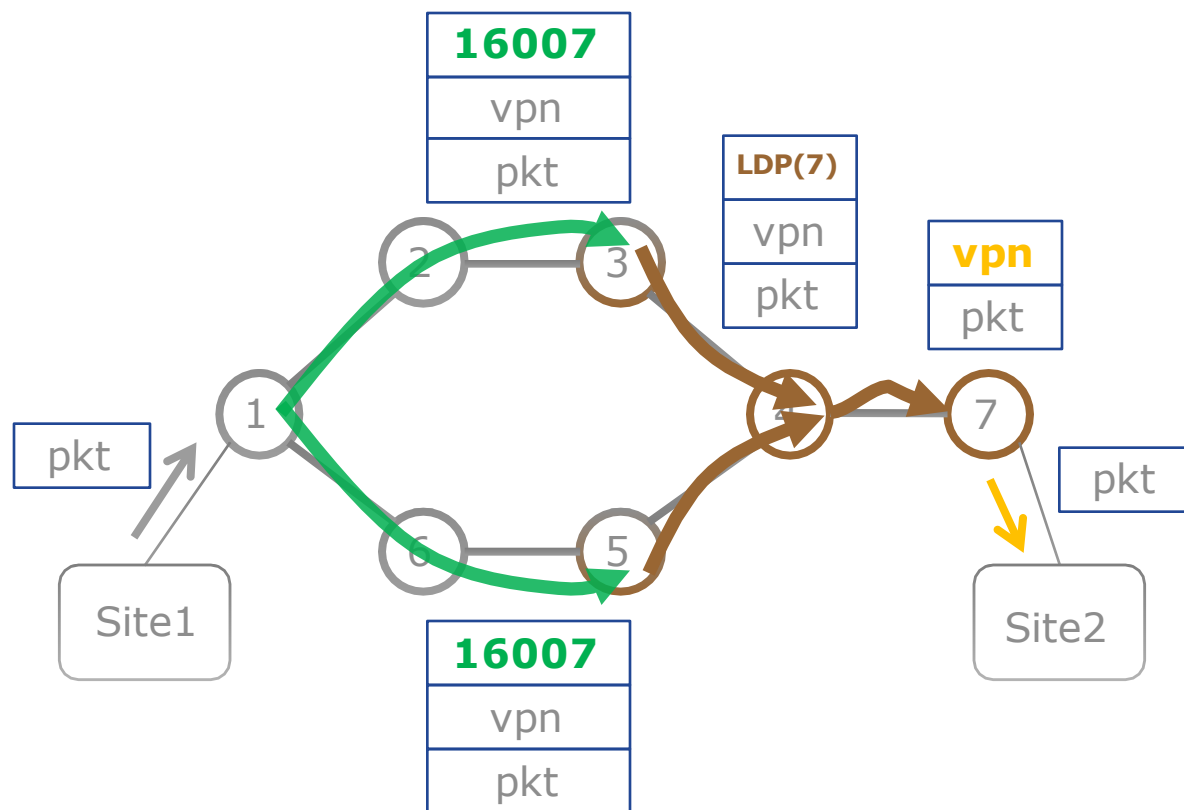
- IGP only
  - No LDP, no RSVP-TE
- ECMP





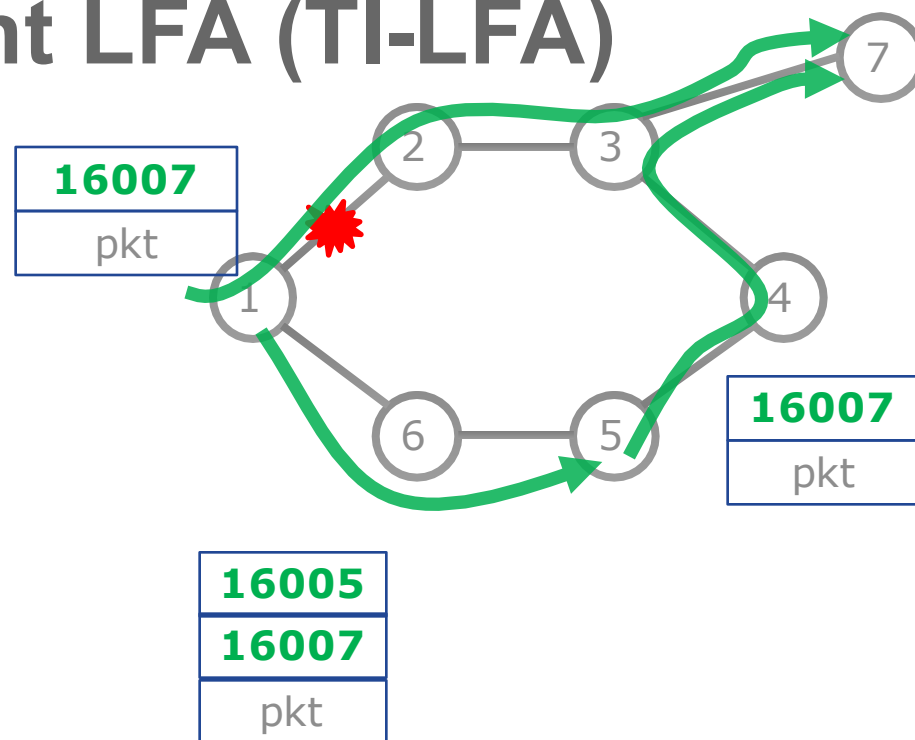
# Seamless interworking with LDP

- Seamless deployment



# Topology-Independent LFA (TI-LFA)

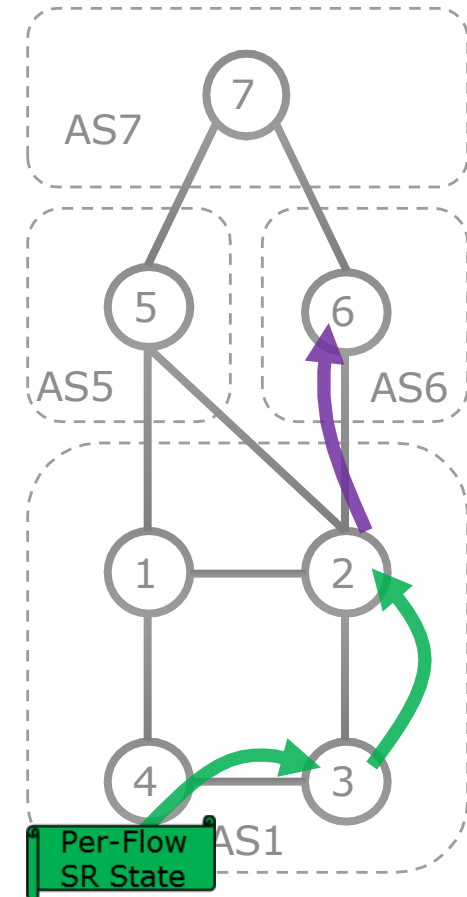
- 50msec FRR in any topology
- IGP Automated
  - No LDP, no RSVP-TE
- Optimum
  - Post-convergence path
- No midpoint backup state
- Detailed operator report
  - S. Litkowski, B. Decraene, Orange
- Mate Design
  - How many backup segments
  - Capacity analysis



# Optimized Content Delivery

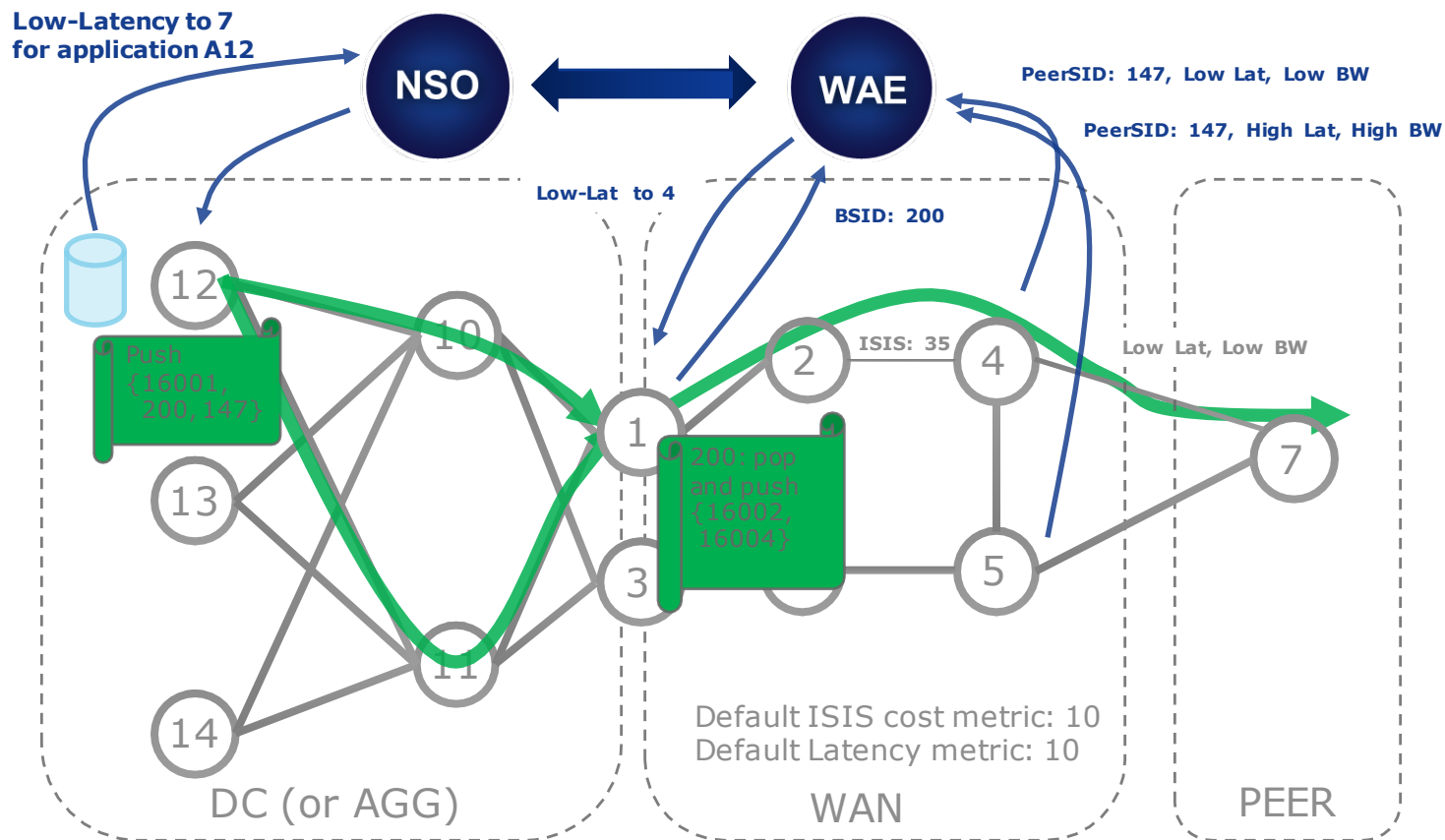
- On a per-content, per-user basis, the content delivery application can engineer
  - the path within the AS
  - the selected border router
  - the selected peer
- Also applicable for engineering egress traffic from DC to peer
  - BGP Prefix and Peering Segments

<b>16003</b>
<b>16002</b>
<b>126</b>
pkt



# Application Engineered Routing

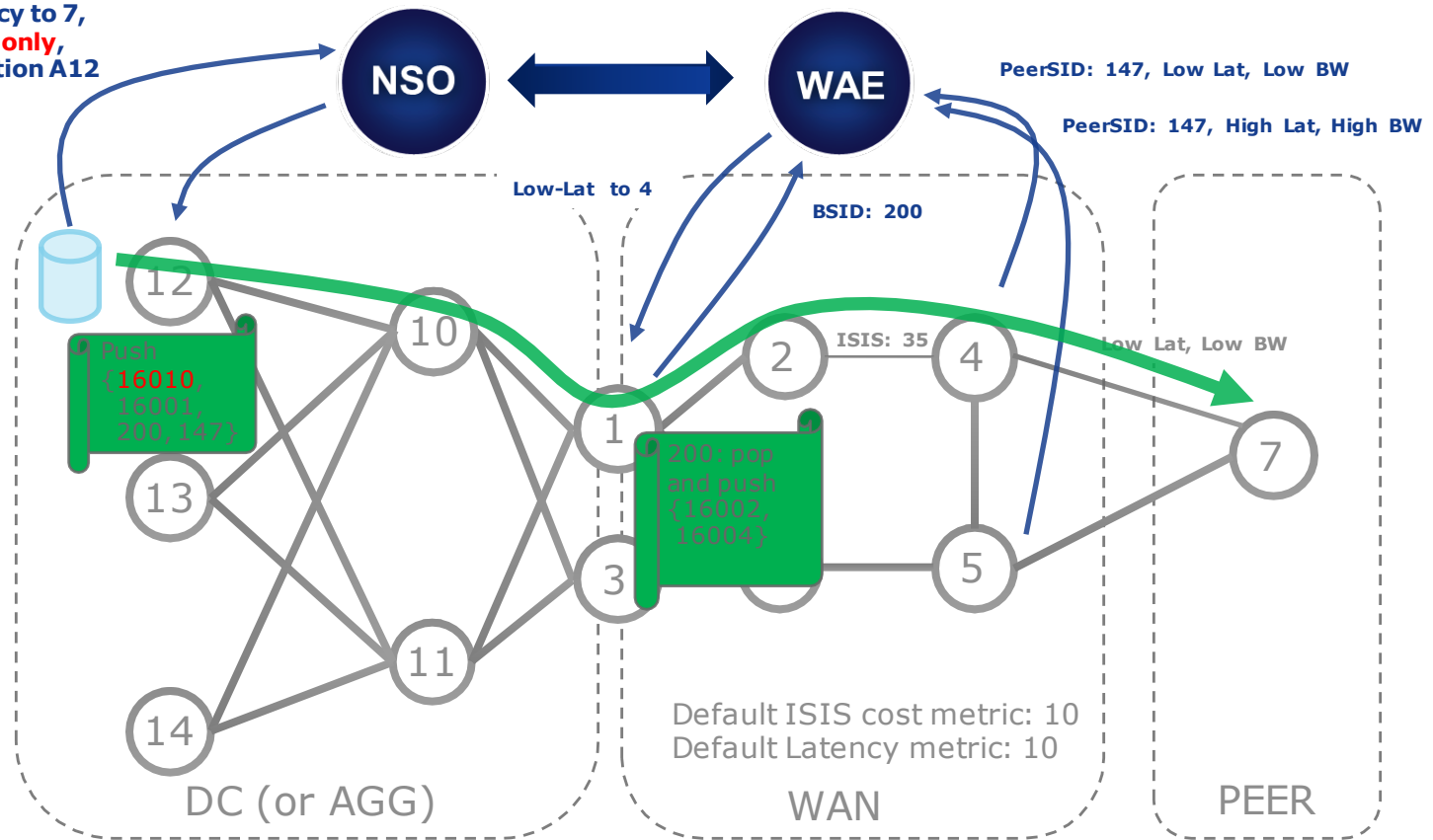
- Per-application flow engineering
  - DC, WAN, AGG, PEER
- End-to-End
  - No signaling
  - No midpoint state
  - No reclassification at boundaries



# Application Engineered Routing

- Per-application flow engineering
  - DC, WAN, AGG, PEER
- End-to-End
  - No signaling
  - No midpoint state
  - No reclassification at boundaries

Low-Latency to 7,  
DC Plane 0 only,  
for application A12

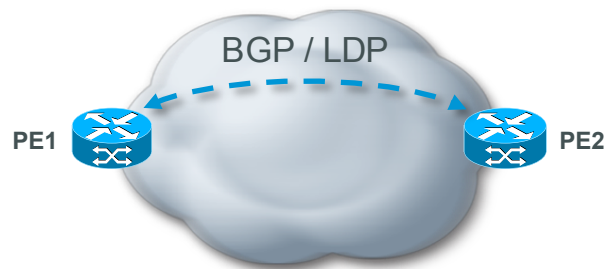




# Control Plane and Data Plane

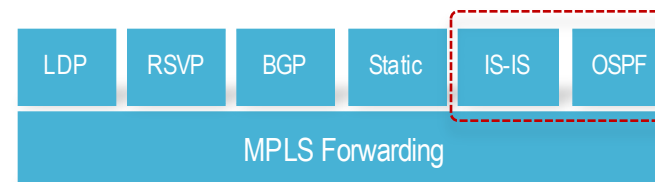
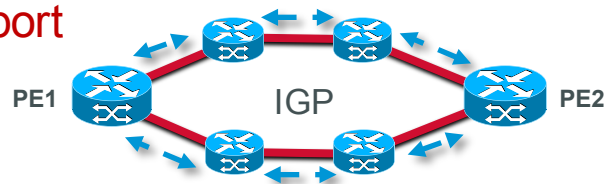
# MPLS Control and Forwarding Operation with Segment Routing

## Services



No changes to control or forwarding plane

## Packet Transport



IGP label distribution for IPv4 and IPv6, same forwarding plane

# SID Encoding

- Prefix SID
  - SID encoded as an index
  - Index represents an offset from SRGB base
  - Index globally unique
  - SRGB may vary across LSRs
  - SRGB (base and range) advertised with router capabilities
- Adjacency SID
  - SID encoded as absolute (i.e. not indexed) value
  - Locally significant
  - Automatically allocated for each adjacency

SR-enabled Node



**SRGB = [ 16000 - 23999 ]**. Advertised as base = 16,000, range = 7,999  
Prefix SID = 16041. Advertised as Prefix SID Index = 41  
Adjacency SID = 24000. Advertised as Adjacency SID = 24000



# SR IS-IS Control Plane Overview

- Level 1, level 2 and multi-level routing
- Prefix Segment ID (Prefix-SID) for host prefixes on loopback interfaces
- Adjacency SIDs for adjacencies
- Prefix-to-SID mapping advertisements (mapping server)
- MPLS penultimate hop popping (PHP) signaling
- MPLS explicit-null label signaling

# IS-IS Configuration

- Required
  - Wide metrics
  - SR enabled under address family IPv4 unicast
- Optional
  - Prefix-SID configured under loopback(s) AF IPv4
- MPLS forwarding enabled automatically on all (non-passive) IS-IS interfaces
- Adjacency-SIDs are automatically allocated for each adjacency

# Configuring Segment Routing for IPv4 Using IS-IS (Cisco IOS XR)

```
router isis DEFAULT
 net 49.0001.1720.1625.5001.00
 address-family ipv4 unicast
  metric-style wide
  segment-routing mpls
 !
 interface Loopback0
  passive
  address-family ipv4 unicast
   prefix-sid absolute 16041
 !
 !
 interface GigabitEthernet0/0/0/0
  point-to-point
  address-family ipv4 unicast
  !
 !
 !
```

Enable Segment Routing for IPv4 with MPLS data plane

Advertise prefix SID 16041 (index 41) for Loopback0

# Configuring Segment Routing for IPv6 Using IS-IS (Cisco IOS XR)

```
router isis DEFAULT
 net 49.0001.1720.1625.5001.00
 address-family ipv6 unicast
  metric-style wide
  segment-routing mpls
 !
 interface Loopback0
  passive
  address-family ipv6 unicast
   prefix-sid absolute 16061
 !
 !
 interface GigabitEthernet0/0/0/0
  point-to-point
  address-family ipv6 unicast
  !
 !
 !
```

Enable Segment Routing for IPv4 with MPLS data plane

Advertise prefix SID 16041 (index 41) for Loopback0

# SR OSPF Control Plane Overview

- IPv4 Prefix Segment ID (Prefix-SID) for host prefixes on loopback interfaces
- MPLS penultimate hop popping (PHP) signaling
- MPLS explicit-null label signaling

# OSPF Configuration

- OSPFv2 control plane
- Required
  - Enable segment-routing under instance or area(s)
    - Command has area scope, usual inheritance applies
  - Enable segment-routing forwarding under instance, area(s) or interface(s)
    - Command has interface scope, usual inheritance applies
- Optional
  - Prefix-SID configured under loopback(s)
- MPLS forwarding enabled on all OSPF interfaces with segment-routing forwarding configured

# Configuring Segment Routing for IPv4 Using OSPF (Cisco IOS XR)

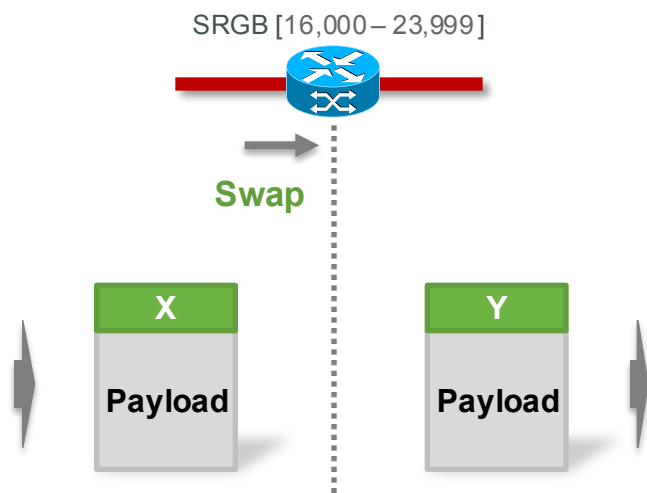
```
router ospf DEFAULT
  router-id 172.16.255.1
  segment-routing mpls
  segment-routing forwarding mpls
  area 0
  interface Loopback0
    passive
    prefix-sid absolute 16041
  !
  interface GigabitEthernet0/0/0/0
    network point-to-point
  !
  !
  !
```

Enable Segment Routing with MPLS data plane

Advertise prefix SID 16041 (index 41) for Loopback0

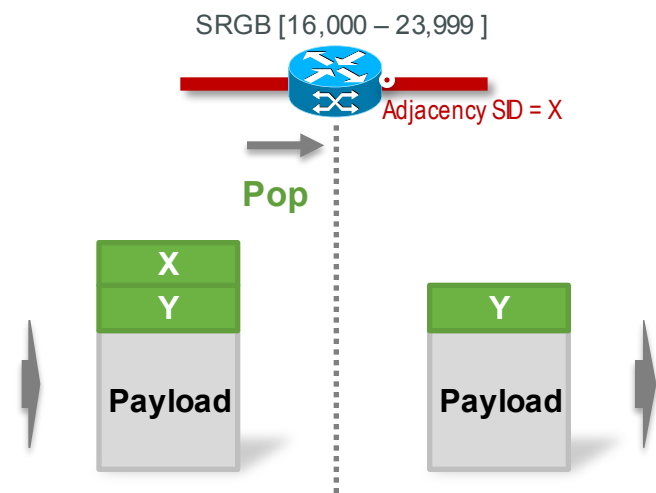
# MPLS Data Plane Operation

## Prefix SID



- Packet forwarded along IGP shortest path
- Packet may leverage ECMP load balancing
- Swap operation performed on input label
- Input label (X) and output label (Y) will have same value when downstream neighbor has same SRGB
- Penultimate hop may perform a pop operation (PHP) if signaled by egress LSR

## Adjacency SID

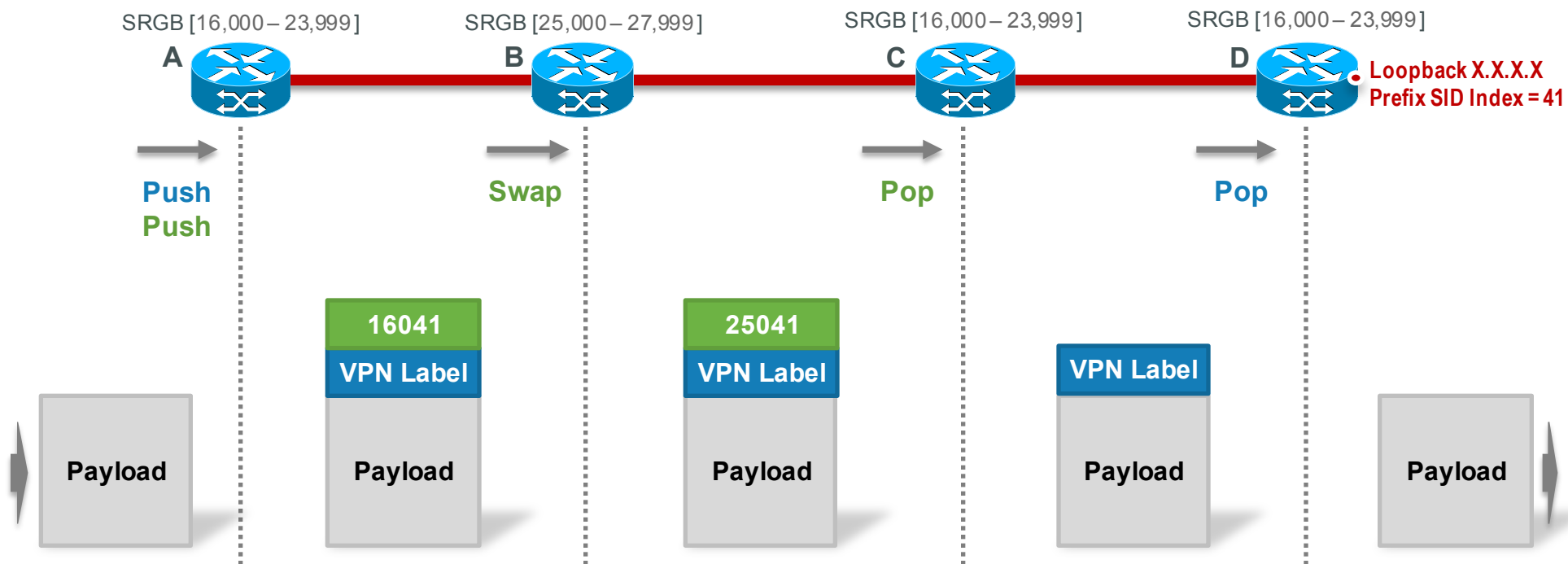


- Packet forwarded along IGP adjacency
- Pop operation performed on input label
- Input topmost label (X) and output label (Y) may or may not have same value
- Penultimate hop always pops last adjacency SID

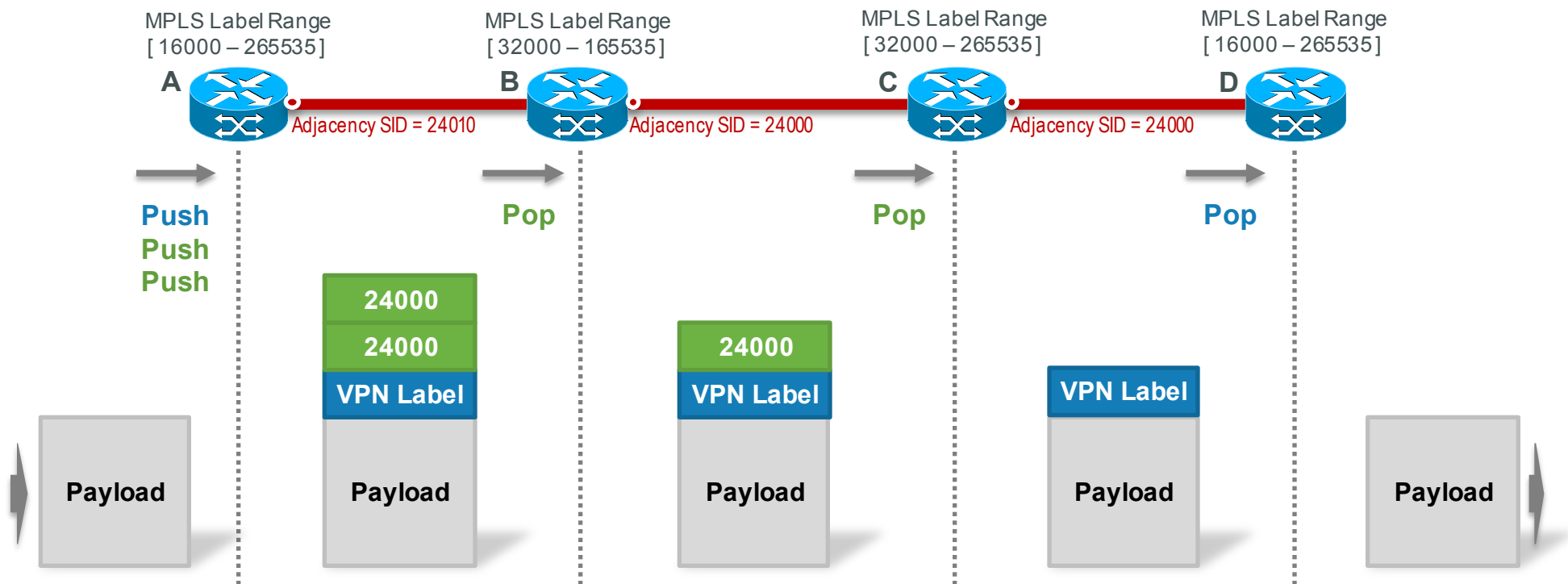




# MPLS Data Plane Operation (Prefix SID)

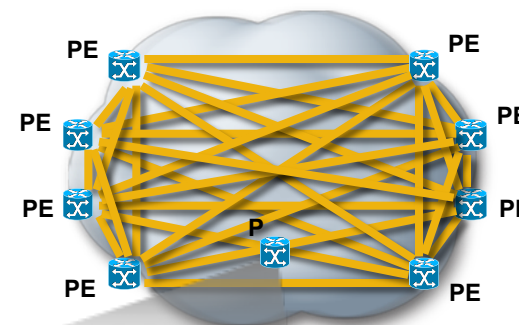


# MPLS Data Plane Operation (Adjacency SIDs)



# MPLS LFIB with Segment Routing

- LFIB populated by IGP (ISIS / OSPF)
- Forwarding table remains constant (Nodes + Adjacencies) regardless of number of paths
- Other protocols (LDP, RSVP, BGP) can still program LFIB



	In Label	Out Label	Out Interface
Network Node Segment Ids	L1	L1	Intf1
	L2	L2	Intf1
	...	...	...
	L8	L8	Intf4
Node Adjacency Segment Ids	L9	L9	Intf2
	L10	Pop	Intf2
	...	...	...
	Ln	Pop	Intf5

**Forwarding table remains constant**



# Traffic Protection

# Topology Independent LFA (TI-LFA) – Benefits

- **100%-coverage** 50-msec link and node protection
- **Simple** to operate and understand
  - automatically computed by the IGP
- **Prevents** transient **congestion** and suboptimal routing
  - leverages the post-convergence path, planned to carry the traffic
- Incremental deployment
  - also **protects LDP traffic**

# Topology Independent LFA – Implementation

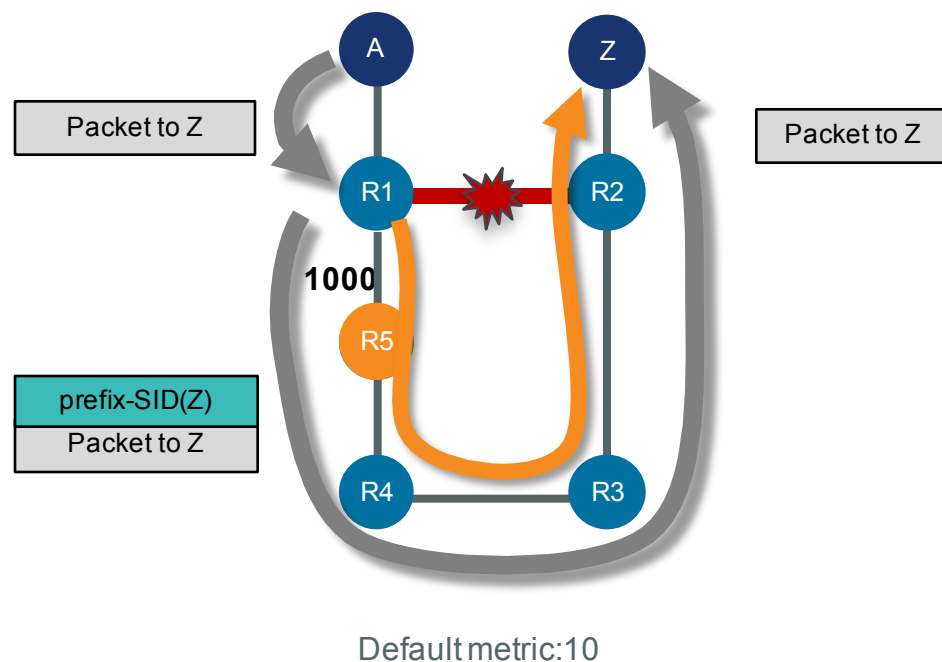
- Leverages existing and proven LFA technology
  - P space: set of nodes reachable from node S (PLR) without using protected link L
  - Q space: set of nodes that can reach destination D without using protected link L
- Enforcing loop-freeness on post-convergence path
  - Where can I release the packet?

At the intersection between the post-convergence shortest path and the Q space
  - How do I reach the release point?

By chaining intermediate segments that are assessed to be loop-free

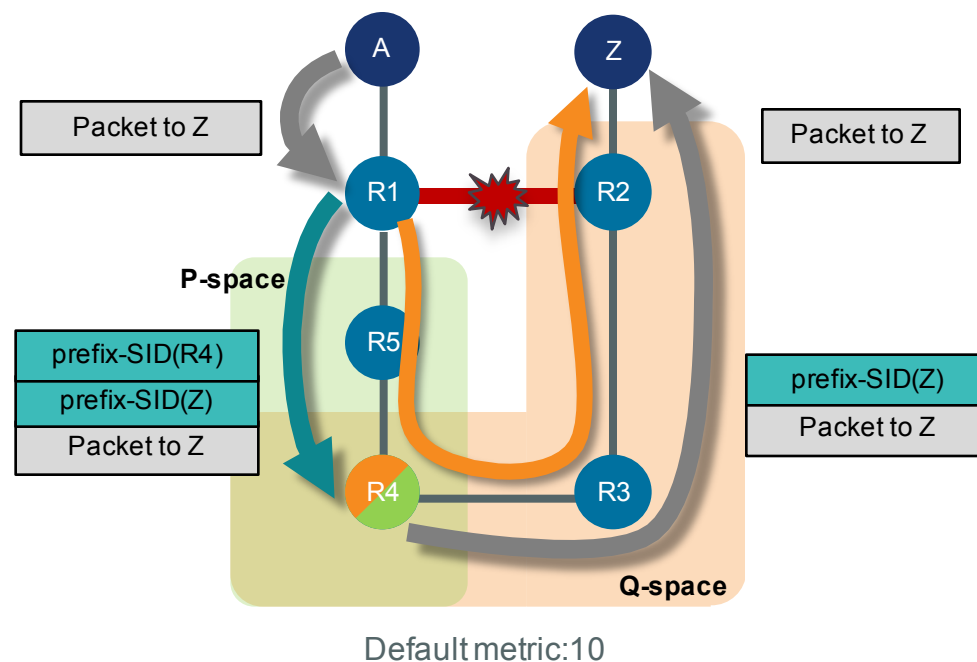
# TI-LFA – Zero-Segment Example

- TI-LFA for link R1R2 on R1
- Calculate LFA(s)
- Calculate post-convergence SPT
- Find LFA on post-convergence SPT
- R1 will steer the traffic towards LFA R5



# TI-LFA – Single-Segment Example

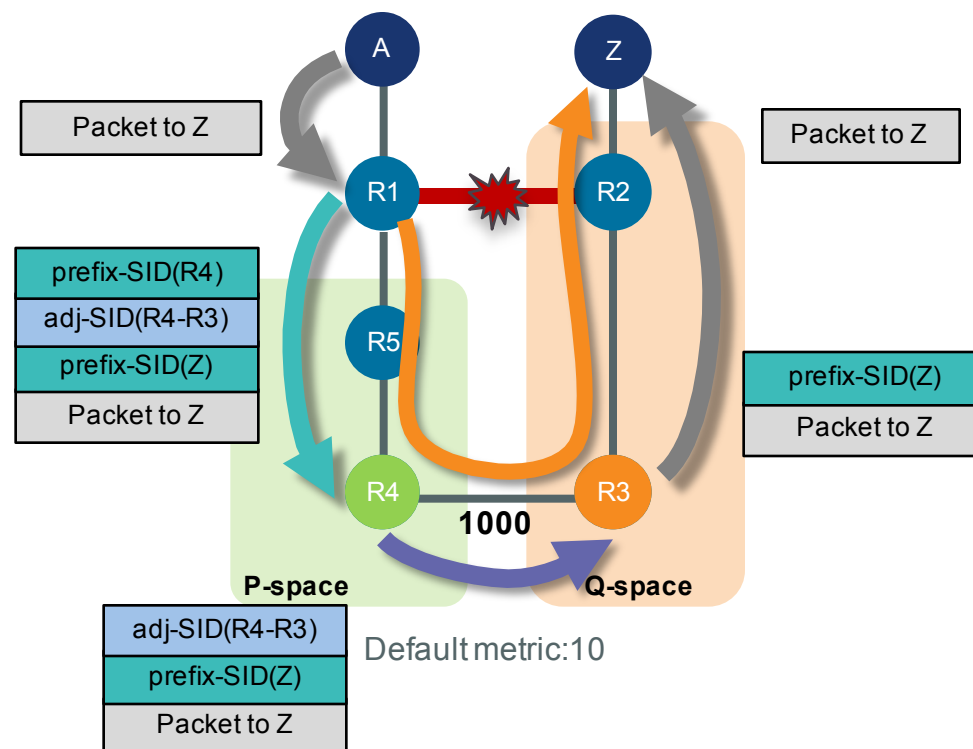
- TI-LFA for link R1R2 on R1
- Calculate P and Q spaces
  - They overlap in this case
- Calculate post-convergence SPT
- Find PQ node on post-convergence SPT
- R1 will push the prefix-SID of R4 on the backup path





# TI-LFA – Double-Segment Example

- TI-LFA for link R1R2 on R1
- Calculate P and Q spaces
- Calculate post-convergence SPT
- Find Q and adjacent P node on post-convergence SPT
- R1 will push the prefix-SID of R4 and the adj-SID of R4-R3 link on the backup path



# Configuring Topology Independent Fast Reroute for IPv4 using Segment Routing and IS-IS (Cisco IOS XR)

```
router isis DEFAULT
  net 49.0001.1720.1625.5001.00
  address-family ipv4 unicast
    metric-style wide
    segment-routing mpls
  !
  interface Loopback0
    passive
    address-family ipv4 unicast
      prefix-sid absolute 16041
    !
  !
  interface GigabitEthernet0/0/0/0
    address-family ipv4 unicast
      fast-reroute per-prefix
      fast-reroute per-prefix ti-lfa
    !
  !
  !
```

Enable TI-LFA for IPv4 prefixes on interface GigabitEthernet0/0/0/0

# Configuring Topology Independent Fast Reroute for IPv6 using Segment Routing and IS-IS (Cisco IOS XR)

```
router isis DEFAULT
 net 49.0001.1720.1625.5001.00
 address-family ipv6 unicast
  metric-style wide
  segment-routing mpls
 !
 interface Loopback0
  passive
  address-family ipv6 unicast
   prefix-sid absolute 16061
 !
 !
 interface GigabitEthernet0/0/0/0
  address-family ipv6 unicast
   fast-reroute per-prefix
   fast-reroute per-prefix ti-lfa
 !
 !
 !
```

Enable TI-LFA for IPv6 prefixes on interface GigabitEthernet0/0/0/0

# Conclusion

- Simple routing extension to enable source routing
- Packet path is determined by prepended segment identifiers (one or more)
- Dataplane agnostic (MPLS, IPv6)
- Network Scalability and agility by reducing network state and simplifying control plane
- Traffic protection with 100% coverage with more optimal routing

Thank you





**CISCO**

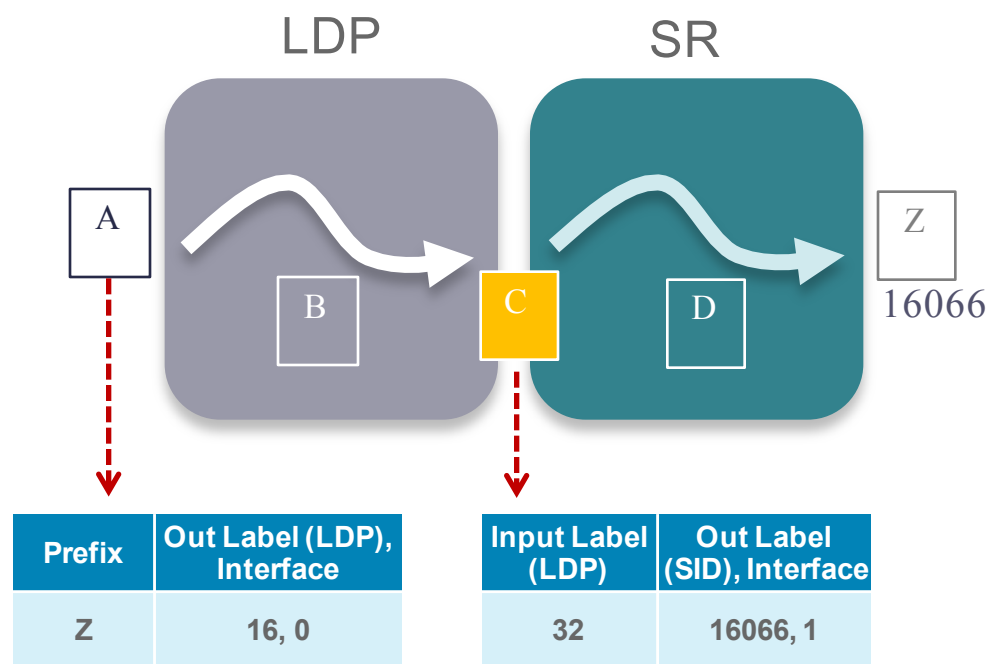
*TOMORROW starts here.*



# SR and LDP Interworking

# LDP to SR

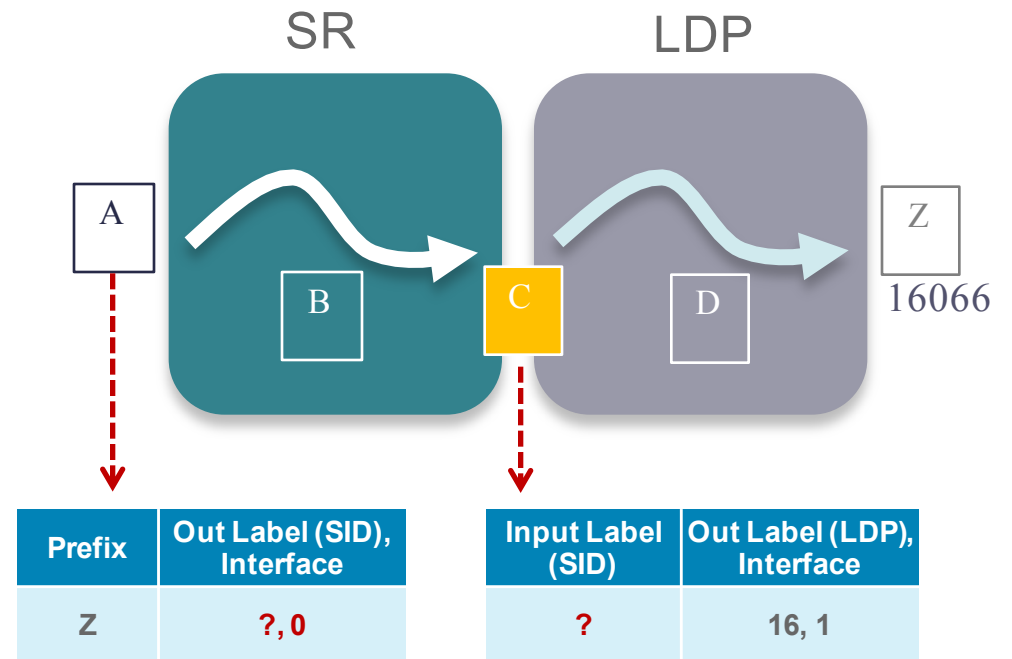
- When a node is LDP capable but its next-hop along the SPT to the destination is not LDP capable
  - no LDP outgoing label
- In this case, the LDP LSP is connected to the prefix segment
- C installs the following LDP-to-SR FIB entry:
  - incoming label: label bound by LDP for FEC Z
  - outgoing label: prefix segment bound to Z
  - outgoing interface: D
- This entry is derived automatically at the routing layer





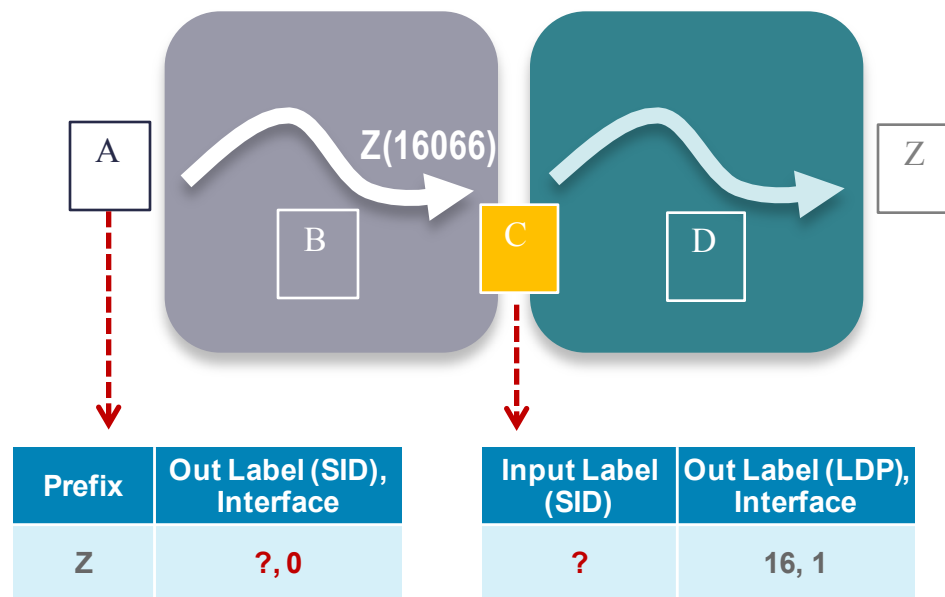
# SR to LDP

- When a node is SR capable but its next-hop along the SPT to the destination is not SR capable
  - no SR outgoing label available
- In this case, the prefix segment is connected to the LDP LSP
  - Any node on the SR/LDP border installs SR-to-LDP FIB entry(ies)



# Mapping Server

- A wants to send traffic to Z, but
  - Z is not SR-capable, Z does not advertise any prefix-SID
 → which label does A have to use?
- The **Mapping Server** advertises the SID mappings for the non-SR routers
  - for example, it advertises that Z is 16068
- A and B install a normal SR prefix segment for 16066
- C realizes that its next hop along the SPT to Z is not SR capable hence C installs an SR-to-LDP FIB entry
  - incoming label: prefix-SID bound to Z (16066)
  - outgoing label: LDP binding from D for FEC Z
- A sends a frame to Z with a single label: 16066



# Active Mapping Policy Preferences

- **Active SID Mapping policy**
  - A set of **non-overlapping** SID mapping entries derived from locally configured SID mappings and SID mappings received from other nodes
- **Backup SID Mapping policy**
  - SID mapping entries that **overlap** with at least one Active SID mapping entry
- When two or more SID mapping entries overlap, which one will be used?
  - Sort all overlapping entries according to preference rules\*
    - Locally configured entries are treated the same as remote entries
  - Only the **most preferred** entry is inserted in the **Active** SID mapping policy
  - The **other** SID-entries are inserted in the **Backup** SID mapping policy

\* Highest router-id > smallest prefix numerical value > smallest first SID value > largest range > latest received

# Configuring a Mapping Server for SR and LDP Interworking for IPv4 Using IS-IS (Cisco IOS XR)

```
router isis DEFAULT
  net 49.0001.1720.1625.5001.00
  address-family ipv4 unicast
  metric-style wide
  segment-routing mpls
  segment-routing prefix-sid-map receive
  segment-routing prefix-sid-map advertise-local
!
...
!
segment-routing
  address-family ipv4
  prefix-sid-map
    172.16.255.1/32 4041 range 8
  !
  !
  !
```

Construct active mapping policy using remotely learned and locally configured mappings (mapping client)

Advertise local mapping policy (mapping server)

Local prefix-to-SID mapping policy  
172.16.255.1/32 - 4041  
:  
172.16.255.8/32 - 4048

# Configuring a Mapping Client for SR and LDP Interworking for IPv4 Using IS-IS (Cisco IOS XR)

```
router isis DEFAULT
  net 49.0001.1720.1625.5001.00
  address-family ipv4 unicast
  metric-style wide
  segment-routing mpls
  segment-routing prefix-sid-map receive
  !
  interface Loopback0
    passive
    address-family ipv4 unicast
      prefix-sid absolute 16041
    !
  !
  interface GigabitEthernet0/0/0/0
    point-to-point
    address-family ipv4 unicast
  !
  !
  !
```

Construct active mapping policy using remotely learned and locally configured mappings (mapping client)